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# Uniform approximation to finite Hilbert transform of oscillatory functions and its algorithm

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## Abstract

For the finite Hilbert transform of oscillatory functions  $Q(f; c, \omega) = \int_{-1}^1 f(x) e^{i\omega x} / (x - c) dt$  with a smooth function  $f$  and real  $\omega \neq 0$ , for  $c \in (-1, 1)$  in the sense of Cauchy principal value or for  $c = \pm 1$  of Hadamard finite-part, we present an approximation method of Clenshaw-Curtis type and its algorithm. Interpolating  $f$  by a polynomial  $p_n$  of degree  $n$  and expanding in terms of the Chebyshev polynomials with  $O(n \log n)$  operations by the FFT, we obtain an approximation  $Q(p_n; c, \omega) \cong Q(f; c, \omega)$ . We write  $Q(p_n; c, \omega)$  as a sum of the sine and cosine integrals and an oscillatory integral of a polynomial of degree  $n - 1$ . We efficiently evaluate the oscillatory integral with a combination of authors' previous method and Keller's method. For  $f(z)$  analytic on the interval  $[-1, 1]$  in the complex plane  $z$ , the error of  $Q(p_n; c, \omega)$  is bounded uniformly with respect to  $c$  and  $\omega$ . Numerical examples illustrate the performance of our method.

*Keywords:* quadrature rule, principal value integral, oscillatory function, Chebyshev interpolation, error analysis, uniform approximation,

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## 1. Introduction

Integrals  $\int_{-1}^1 f(t)K(t) dt$  of a real-valued, regular function  $f(t)$  and the singular or/and oscillatory kernel  $K(t) = 1/(t - c)$ ,  $e^{i\omega t}$  ( $i = \sqrt{-1}$ , real  $\omega \neq 0$ ) or

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$e^{i\omega t}/(t-c)$ , often arise in scientific computing (cf. [1, 1.6, 2.10]).

In this work, we consider integrals of  $K(t) = e^{i\omega t}/(t-c)$  ( $c \in [-1, 1]$ ), finite Hilbert transform of highly-oscillatory functions (assuming  $|\omega| \gg 1$ )

$$Q(f; c, \omega) := \begin{cases} \mathcal{P} \int_{-1}^1 \frac{f(t) e^{i\omega t}}{t-c} dt, & c \in (-1, 1), \\ \mathcal{F} \int_{-1}^1 \frac{f(t) e^{i\omega t}}{t-c} dt, & c = \pm 1, \end{cases} \quad (1.1)$$

5 in the sense of Cauchy principal value (CPV) (cf. [1, p.11]) for  $c \in (-1, 1)$ , or of Hadamard finite-part (HFP) (cf. [1, 1.6.1]) for  $c = \pm 1$ . High oscillation and singularity make it difficult to approximate the integrals with well-known quadrature rules such as the Gaussian and Newton-Cotes rules.

For efficiently approximating  $Q(f; c, \omega)$  (1.1), we develop a method, an extended Clenshaw–Curtis (C–C) quadrature rule [2] (cf. [3, Sec.8.3]). We use approximation methods [4–6] for CPV integrals and an algorithm which is a combination of two methods given in [7, 8] for evaluating oscillatory integrals (cf. [9, 10]). For a C–C type method for CPV integrals with logarithmic singularity, see [11]. Since the C–C rule has some advantages shown below, several  
15 C–C type methods for  $Q(f; c, \omega)$  have been developed (cf. [12–14]). We review them later. For some works other than the C–C type method, see [15–17].

The C–C rule is an interpolatory integration rule. For an integer  $n \geq 1$ , let  $p_n$  be a polynomial interpolating  $f$  at the Chebyshev points  $x_{n,j} = \cos(\pi j/n)$  ( $0 \leq j \leq n$ ) on  $[-1, 1]$ . Then, we obtain an approximation  $Q(p_n; c, \omega) \cong Q(f; c, \omega)$  in (1.1). We expand  $p_n$  in terms of the Chebyshev polynomials of the first kind  $T_k(t) = \cos k\theta$  ( $t = \cos \theta$ ),

$$f(t) \cong p_n(t) = \sum_{k=0}^n{}'' a_k T_k(t), \quad (1.2)$$

where the double prime indicates the summation whose first and last terms are halved. The coefficients  $a_k$  are efficiently evaluated with  $O(n \log_2 n)$  arithmetic operations with the fast Fourier transform (FFT) (cf. [18],[19, p.232]). If  $f(t)$  is  
20 a sufficiently smooth function, then  $p_n$  in (1.2) converges rapidly to  $f$  as  $n$  grows. We can effectively increase the accuracy of the approximation by doubling the

number of the nodes  $\{x_{n,j}\}$  ( $0 \leq j \leq n$ ) for  $p_n$ . Indeed, these nodes are reused in  $\{x_{2n,j}\}$  ( $0 \leq j \leq 2n$ ) for  $p_{2n}$ , since  $\{x_{2n,j}\}_{j=0}^{2n} = \{x_{n,j}\}_{j=0}^n \cup \{\cos[\pi(2j+1)/(2n)]\}_{j=0}^{n-1}$ ; see [20–22], [23, p.75].

25 To evaluate the approximation  $Q(p_n; c, \omega)$ , every method of C–C type in [12–14] makes use of a *recurrence relation* (RR), although a numerical instability often occurs in the computation of the RR (cf. [24, Section 6]). We review them and ours.

Keller [13] approximates the indefinite integral,

$$\int \frac{(p_n(t) - R)e^{i\omega t}}{t - c} dt \cong e^{i\omega t} \sum_{k=0}^N {}' \alpha_k T_k(t) \quad (-1 \leq t \leq 1), \quad (1.3)$$

where  $N = n(1 + o(1))$  and  $R \cong p_n(c)$ . The prime indicates the summation whose first term is halved. The coefficients  $\alpha_k$  satisfy the five-term RR,

$$\begin{aligned} \frac{i\omega}{2} \alpha_{k-2} + (k-1 - i\omega c) \alpha_{k-1} - 2kc \alpha_k + (k+1 + i\omega c) \alpha_{k+1} - \frac{i\omega}{2} \alpha_{k+2} \\ = a_{k-1} - a_{k+1} \quad (k \geq 2). \end{aligned} \quad (1.4)$$

Then,  $Q(p_n; c, \omega)$  for  $c \in (-1, 1)$  is given by

$$Q(p_n; c, \omega) \cong e^{i\omega} \sum_{k=0}^N {}' \alpha_k - e^{-i\omega} \sum_{k=0}^N {}' (-1)^k \alpha_k + R \int_{-1}^1 \frac{e^{i\omega t}}{t - c} dt, \quad (1.5)$$

30 where  $R = (a_0 - a_2)/2 - i\omega c(\alpha_2 - \alpha_0)/2 - i\omega(\alpha_1 - \alpha_3)/4 - \alpha_2 + c\alpha_1$ . The integrals on the right of (1.5) is simple to evaluate, since they are expressed by the sine and cosine integrals or the exponential integral; see Section 2 and [12–14].

Other works in [12, 14] similarly and exactly divide  $Q(p_n; c, \omega)$  into a non-singular integral  $I(p_n; c, \omega)$  and CPV or HFP integrals,

$$Q(p_n; c, \omega) = I(p_n; c, \omega) + p_n(c) \begin{cases} \int_{-1}^1 \frac{e^{i\omega t}}{t - c} dt & (c \in (-1, 1)), \\ \int_{-1}^1 \frac{e^{i\omega t}}{t - c} dt & (c = \pm 1), \end{cases} \quad (1.6)$$

$$I(p_n; c, \omega) := \int_{-1}^1 \frac{p_n(t) - p_n(c)}{t - c} e^{i\omega t} dt. \quad (1.7)$$

Clenshaw's algorithm effectively evaluate  $p_n(c)$  at  $t = c$  (cf. [3, p.27]); see (5.1). The problem is to efficiently evaluate the integral  $I(p_n; c, \omega)$  (1.7).

Wang, Zhang and Huybrechs [14], an improvement of the method by Wang and Xiang [12], make use of the key relation (cf. [4])

$$\frac{p_n(t) - p_n(c)}{t - c} = \sum_{k=0}^{n-1} b_k T_k(t) =: q_{n-1}(t), \quad (1.8)$$

to obtain a simple form of the integral in (1.7),

$$I(p_n; c, \omega) = \int_{-1}^1 q_{n-1}(t) e^{i\omega t} dt = \sum_{k=0}^{n-1} b_k \nu_k, \quad \nu_k := \int_{-1}^1 T_k(t) e^{i\omega t} dt. \quad (1.9)$$

The coefficients  $b_k$  in (1.8) are obtained by backward computation of the three-term RR (3.2) below. The moments  $\nu_k$  in (1.9) are given by

$$\nu_k = \frac{1}{i\omega} \{e^{i\omega} - (-1)^k e^{-i\omega}\} - \frac{k}{i\omega} \mu_{k-1} \quad (k \geq 1),$$

where  $\mu_k = \int_{-1}^1 U_k(t) e^{i\omega t} dt$ , and  $U_k(t) = \sin(k+1)\theta / \sin \theta$  ( $t = \cos \theta$ ) is the Chebyshev polynomial of the second kind of degree  $k$  (cf. [3, p.3]). The moments  $\mu_k$  satisfy the three-term RR,

$$\mu_k + \frac{2k}{i\omega} \mu_{k-1} - \mu_{k-2} = \frac{2}{i\omega} \{e^{i\omega} - (-1)^k e^{-i\omega}\} \quad (k \geq 2). \quad (1.10)$$

To evaluate  $I(p_n; c, \omega)$  in (1.9), we show our previous method [8] now. Later we give an improved algorithm with a smaller number of arithmetic operations. We expand the indefinite integral in terms of the Chebyshev polynomials,

$$\int q_{n-1}(t) e^{i\omega t} dt = \frac{e^{i\omega t}}{i\omega} \varphi(t), \quad \varphi(t) := \sum_{k=0}^{\infty} d_k T_k(t). \quad (1.11)$$

Then, we obtain

$$I(p_n; c, \omega) = \frac{e^{i\omega} \varphi(1) - e^{-i\omega} \varphi(-1)}{i\omega} = \frac{e^{i\omega}}{i\omega} \sum_{k=0}^{\infty} d_k - \frac{e^{-i\omega}}{i\omega} \sum_{k=0}^{\infty} (-1)^k d_k. \quad (1.12)$$

Since  $\varphi(t)$  in (1.11) is an entire function (cf. [8]), the coefficients  $d_k$  converge very fast to 0 as  $k \rightarrow \infty$ . They satisfy the three-term RR,

$$d_{k-1} + \frac{2k}{i\omega} d_k - d_{k+1} = b_{k-1} - b_{k+1} \quad (k = 1, 2, \dots), \quad (1.13)$$

with the coefficients  $b_k$  of  $q_{n-1}(t)$  in (1.8) and  $b_k = 0$  ( $k \geq n$ ); see [8, Theorem 2.1]. Note that the left-hand sides of (1.10) and (1.13) are almost the same.

The numerical instability of the RRs (1.4), (1.10) and (1.13) is a serious problem. See [24, Section 8] and [25, Section 3.3] for the stability of RRs of an arbitrary order. If  $n - 1 \leq |\omega|$ , then forward and backward computations of the RRs are numerically stable (cf. [26]). But, when  $n - 1 > |\omega|$ , neither  
40 forward nor backward computation of the RRs is stable owing to severe loss of significant figures occurring in the computation process. To avoid this difficulty it is required to transform the RR into a system of linear equations of size  $N \geq n - 1$ , as shown below. It is desirable to determine  $N$  as small as possible. Unique solution of the linear system requires some additional boundary conditions (cf.  
45 [25, 27–29]).

For the linear system (1.4), Keller [13] uses Lozier’s algorithm [25] with one initial condition  $\alpha_{\lfloor |\omega| \rfloor} = 0$  and three trailing conditions  $\alpha_{N+1} = \alpha_{N+2} = \alpha_{N+3} = 0$  for a properly selected value of  $N$ . By assuming the asymptotic convergence rate of  $|\alpha_k|$  as  $k \rightarrow \infty$ , the value of  $N$  is determined so that

$$\max\{|\alpha_{N-3}|, |\alpha_{N-2}|, |\alpha_{N-1}|, |\alpha_N|\} < \epsilon \max_{-1 \leq x \leq 1} |p_n(x)|/N^2, \quad (1.14)$$

with a machine epsilon  $\epsilon$ . To estimate the smallest value of  $N$  is not easy since it requires to estimate the magnitude of  $|\alpha_k|$  accurately and efficiently. The URL to Matlab codes implementing the method is shown in [13].

Wang, Zhang and Huybrechs [14] adopt the method due to Domínguez et al. [29] to solve the linear system (1.10) for  $N \geq n - 1$ . In addition to the  
50 trivial boundary condition  $\mu_0 = (2 \sin \omega)/\omega$ , a sufficiently accurate value of  $\mu_N$  is needed. Using an asymptotic expansion of  $\mu_N$ , Domínguez et al. devise an ingenious method to obtain the approximate value  $\tilde{\mu}_N$  of  $\mu_N$  for a sufficiently large  $N$ . Again, to determine the smallest value of  $N$  is not easy since it  
55 requires to estimate the error of  $\tilde{\mu}_N$  precisely and efficiently. Moreover, in order to guarantee the accuracy of  $\tilde{\mu}_N$ , the value of  $N$  can not be so small. An algorithm implementing this method is given (cf. [14]).

Finally, we outline our scheme to solve the linear system (1.13). The normalization relation  $\varphi(-1) = \sum_{k=0}^{\infty} (-1)^k d_k = 0$  is used in [8] for a boundary condition. Here, we choose the condition  $d_{\lfloor |\omega| \rfloor} = 0$  given in [7] that leads to a

simpler and faster algorithm. See [7] and Subsection 3.2 and Section 4 for details and a comparison of the computational costs. Since  $d_k$  converges rapidly to 0 ( $k \rightarrow \infty$ ), instead of the infinite summation (1.11) for  $\varphi(t)$ , it is sufficient to use a truncated Chebyshev series

$$\varphi^{[N]}(t) = \sum_{k=0}^N d_k^{[N]} T_k(t) \approx \varphi(t), \quad (1.15)$$

with an integer  $N \geq n - 1$ . Then, in view of (1.15) we have an approximation  $I^{[N]}(p_n; c, \omega)$  to  $I(p_n; c, \omega)$  in (1.12),

$$I^{[N]}(p_n; c, \omega) = \frac{e^{i\omega}}{i\omega} \sum_{k=0}^N d_k^{[N]} - \frac{e^{-i\omega}}{i\omega} \sum_{k=0}^N (-1)^k d_k^{[N]}. \quad (1.16)$$

As shown in Lemma 1.1 (cf. [8, Theorem 2.3]) below, we easily and automatically determine the value of  $N$  close to  $n - 1$  such that the error of  $I^{[N]}(p_n; c, \omega)$  is  
60 at the level of the unit roundoff  $u$  of computer (cf. [30, p.38]).

**Lemma 1.1.** *The truncation error  $E_T$  of the approximation  $I^{[N]}(p_n; c, \omega)$  in (1.16), where  $N \geq n - 1$ , is bounded by the last coefficient  $d_N^{[N]}$  of  $\varphi^{[N]}(x)$  in (1.15) as*

$$E_T := |I(p_n; c, \omega) - I^{[N]}(p_n; c, \omega)| \leq |d_N^{[N]}|. \quad (1.17)$$

Lemma 1.1 is proved almost in the same way as that in [8]; see Appendix A.

We solve the system of linear equations by Gaussian elimination with  $O(N)$  arithmetic operations. We increase  $N$  until  $|d_N^{[N]}| \leq u$  is satisfied in the forward elimination process, followed by computing  $d_k^{[N]}$  ( $0 \leq k \leq N - 1$ ) by back  
65 substitution. The result guarantees that  $E_T \leq u$ , from Lemma 1.1.

Lemma 1.1 provides an efficient and verified error bound to determine  $N$ . As we have mentioned above, on the other hand, it is not easy to determine the smallest value of  $N$  in [13, 14]. In the scheme of Domínguez et al. [29] that is of a different type from ours, the size  $N$  of the linear system  
70 is fixed, provided a sufficiently accurate approximation  $\tilde{\mu}_N$  to  $\mu_N$  is given. To obtain  $\tilde{\mu}_N$ , however, the value of  $N$  is doubled from  $n - 1$  until the required accuracy is achieved. Using the Matlab routines by Domínguez et al.,

Table 1: The size  $N$  of linear systems for each  $n - 1$

$n - 1 =$	80	160	320	640
Domínguez et al.	320	320	640	1280
Ours	<b>110</b>	<b>173</b>	<b>322</b>	<b>640</b>

we performed numerical experiments for the integral  $\int_{-1}^1 f(x)e^{i\omega x} dx$ , where  $f(x) = (1 - \alpha)/(1 - 2\alpha x + \alpha^2) = 2 \sum_{k=0}^{\infty} \alpha^k T_k(x)$  with  $\alpha = 0.9$ . We compare  
75 the size  $N$  of our linear system and  $N$  in [29] required to within the unit round-off,  $u \cong 1.11 \times 10^{-16}$  of computer. When  $\omega = 60$  and  $n - 1 = 80, 160, 320$  and 640, the required values of  $N$  in Domínguez et al. [29] and our method are given in Table 1. The Matlab code by Keller [13], where a strange formula different from (1.14) is used, indicates that  $N = 117$  ( $n - 1 \leq 115$ ) and  $N = n + 1$   
80 ( $n - 1 > 115$ ) when  $\omega = 60$ .

Our method approximates  $Q(f; c, \omega)$  uniformly, independently of  $c$  and  $\omega$ ; see Theorem 6.5, and is efficiently implemented.

The paper is organized as follows. Section 2 shows that the integrals in (1.6) are expressed by the sine and cosine integrals. Section 3 gives a review of  
85 [8] for approximating the oscillatory integrals  $I(p_n; c, \omega)$  and its improvement. In Section 4 we rewrite the RR (1.13) in the form of a linear system of finite dimension and solve it by Gaussian elimination via bordering (cf. [31, p.55]). The number of arithmetic operations is examined. Section 5 summarizes the algorithm to compute  $Q(p_n; c, \omega)$ . In Section 6, assuming that  $f(z)$  is analytic  
90 on  $[-1, 1]$  in the complex plane  $z$ , we give error analysis of  $Q(p_n; c, \omega)$  and a uniform error bound with respect to  $c$  and  $\omega$ . Section 7 gives numerical examples to show the error behavior of  $Q(p_n; c, \omega)$ .



## 2. Treatment of CPV and HFP integrals

The CPV and HFP integrals in (1.6) are expressed by the sine and cosine  
 95 integrals or the exponential integral (cf. [12–16]).

**Lemma 2.1.** *The CPV integral  $\oint_{-1}^1 e^{i\omega t}/(t-c) dt$  with  $c \in (-1, 1)$  in (1.6) is written by*

$$\begin{aligned} \oint_{-1}^1 \frac{e^{i\omega t}}{t-c} dt &= e^{i\omega c} [\text{Ci}(|(1-c)\omega|) - \text{Ci}(|(1+c)\omega|)] \\ &\quad + i\{\text{Si}((1-c)\omega) + \text{Si}((1+c)\omega)\}, \quad c \in (-1, 1), \end{aligned} \quad (2.1)$$

where  $\text{Si}(x)$  and  $\text{Ci}(x)$  are the sine and cosine integrals (cf. [32, 5.2.1, 5.2.27], [33, p.181, p.187]), respectively, and given by

$$\text{Si}(x) = \int_0^x \frac{\sin t}{t} dt, \quad \text{Ci}(x) = \begin{cases} -\int_x^\infty \frac{\cos t}{t} dt & (x > 0), \\ \pi i - \int_x^\infty \frac{\cos t}{t} dt & (x < 0). \end{cases} \quad (2.2)$$

See [12, 15] for the proof.

**Lemma 2.2.** *For the HFP integral  $\oint_{-1}^1 e^{i\omega t}/(t-c) dt$  with  $c = \pm 1$  in (1.6), we have*

$$\oint_{-1}^1 \frac{e^{i\omega t}}{t-c} dt = \begin{cases} e^{i\omega c} [c\{\gamma + \log|\omega| - \text{Ci}(2|\omega|)\} + i\text{Si}(2\omega)] & (\omega \neq 0), \\ -c \log 2 & (\omega = 0), \end{cases} \quad (2.3)$$

with Euler's constant  $\gamma = 0.5772 \dots$ .

The formula similar to (2.3) given in terms of the exponential integral is an immediate consequence of the results presented in [14].

*Proof of Lemma 2.2* The case  $\omega = 0$  is trivial. For  $\omega \neq 0$  we claim that

$$\oint_{-1}^1 \frac{e^{i\omega(t-c)} - 1}{t-c} dt = c\{\gamma + \log(2|\omega|) - \text{Ci}(2|\omega|)\} + i\text{Si}(2\omega). \quad (2.4)$$

Then, we verify (2.3) for  $\omega \neq 0$  using (2.4) and  $\oint_{-1}^1 1/(t-c) dt = -c \log 2$  in the relation

$$\oint_{-1}^1 \frac{e^{i\omega t}}{t-c} dt = e^{i\omega c} \left\{ \int_{-1}^1 \frac{e^{i\omega(t-c)} - 1}{t-c} dt + \oint_{-1}^1 \frac{1}{t-c} dt \right\}.$$

It remains to verify (2.4). For  $c = \pm 1$ , by the change of variables  $u = c\omega(c - t) = \omega(1 - ct)$  we have

$$\begin{aligned} \int_{-1}^1 \frac{e^{i\omega(t-c)} - 1}{t - c} dt &= \int_{(1+c)\omega}^{(1-c)\omega} \frac{e^{-icu} - 1}{u} du = -c \int_0^{2\omega} \frac{e^{-icu} - 1}{u} du \\ &= c \int_0^{2\omega} \frac{1 - \cos u}{u} du + i \int_0^{2\omega} \frac{\sin u}{u} du. \end{aligned} \quad (2.5)$$

Abramowitz & Stegun [32, 5.2.2] gives

$$\int_0^x \frac{1 - \cos u}{u} du = \gamma + \log x - \text{Ci}(x) \quad (x > 0).$$

Since the left-hand side above is an even function, it follows that

$$\int_0^x \frac{1 - \cos u}{u} du = \gamma + \log |x| - \text{Ci}(|x|) \quad (x \neq 0). \quad (2.6)$$

100 Using (2.2) and (2.6) in (2.5) verifies (2.4).  $\square$

### 3. Evaluation of $a_k$ , $b_k$ and $I(p_n; c, \omega)$

We give an expression for the Chebyshev coefficients  $a_k$  in (1.2) and  $b_k$  in (1.8) followed by the evaluation of  $I(p_n; c, \omega)$ .

#### 3.1. Evaluation of $a_k$ and $b_k$

The interpolation condition  $f(x_{n,j}) = p_n(x_{n,j})$  ( $0 \leq j \leq n$ ) gives the coefficients  $a_k$  in (1.2) (cf. [3, p.156]),

$$a_k = \frac{2}{n} \sum_{j=0}^n f\left(\cos \frac{\pi j}{n}\right) \cos \frac{\pi j k}{n} \quad (0 \leq k \leq n). \quad (3.1)$$

105 The summations on the right are evaluated with  $O(n \log_2 n)$  multiplications with the FFT (cf. [18]); see Section 5. If  $f$  is sufficiently smooth, then  $|a_k|$  decreases rapidly to zero as  $k$  grows, therefore  $p_n$  in (1.2) converges rapidly to  $f$  as  $n$  grows (cf. [34, p.175], [21]).

The coefficients  $b_k$  for  $q_{n-1}$  in (1.8) satisfy the three-term RR (cf. [4]),

$$b_{k-1} - 2c b_k + b_{k+1} = 2a_k \quad (k = 1, 2, \dots, n), \quad (3.2)$$

with  $a_n$  replaced by  $a_n/2$ . We set  $b_k = 0$  ( $k \geq n$ ) for convenience. Since  
110  $c \in [-1, 1]$ , backward computation of the RR (3.2) is numerically stable and gives the values of  $b_0 \dots, b_{n-1}$  with starting values  $b_n = b_{n+1} = 0$ .

### 3.2. Evaluation of $I(p_n; c, \omega)$

Differentiating both the sides of the first equation in (1.11) gives the linear first-order differential equation for  $\varphi(t)$ ,

$$\varphi'(t) + i\omega \varphi(t) = i\omega q_{n-1}(t) \quad (-1 \leq t \leq 1), \quad (3.3)$$

with  $q_{n-1}(t)$  in (1.8). Lemma 3.1 gives the general solution of (3.3).

**Lemma 3.1** (cf. [8]). *The general solution  $\varphi(t)$  of the differential equation (3.3) is an entire function given, with a constant  $C$ , by*

$$\varphi(t) = Ce^{-i\omega t} + \phi_{n-1}(t), \quad \phi_{n-1}(t) := \sum_{k=0}^{n-1} \left(\frac{i}{\omega}\right)^k \frac{d^k}{dt^k} q_{n-1}(t). \quad (3.4)$$

We expand the polynomial  $\phi_{n-1}(t)$  of degree  $n-1$  and  $e^{-i\omega t}$  in (3.4) in terms of the Chebyshev polynomials,

$$\phi_{n-1}(t) = \sum_{k=0}^{n-1} d_k^{(0)} T_k(t), \quad (3.5)$$

$$e^{-i\omega t} = 2 \sum_{k=0}^{\infty} (-i)^k J_k(\omega) T_k(t), \quad (3.6)$$

with the Bessel function of the first kind  $J_k(z)$  (cf. [3, p.109], [32, p.358, p.375]). Inserting (3.5) and (3.6) into (3.4) for  $\varphi(t)$  and comparing  $d_k$  for  $\varphi(t)$  in (1.11), we obtain

$$d_k = \begin{cases} 2C(-i)^k J_k(\omega) + d_k^{(0)} & (0 \leq k \leq n-1), \\ 2C(-i)^k J_k(\omega) & (n \leq k). \end{cases} \quad (3.7)$$

Lemma 3.1 implies two types of the solution  $\varphi(t)$  with  $C \neq 0$  or  $C = 0$ .  
 115 Since the definite integral  $I(p_n; c, \omega)$  defined by (1.7) is independent of  $C$ , an arbitrary value of  $C$  is possible provided  $|C| < \infty$ . When  $C \neq 0$ , the function  $\varphi(t)$  behaves more like  $e^{-i\omega t}$  and the Chebyshev expansion of  $\varphi(t)$  in (1.11) converges very fast. This is seen as follows. Since  $J_k(z) \sim \{ez/(2k)\}^k / \sqrt{2\pi k}$  ( $k \rightarrow \infty$ ) (cf. [32, p.365]), it follows that  $J_k(\omega)$ , and, consequently,  $d_k$  in (3.7)  
 120 converge to 0 very fast for  $k > |\omega|$  with  $\omega$  fixed.

Let  $M = \lfloor |\omega| \rfloor$ . When  $n-1 \leq M$ , we use the polynomial solution with  $C = 0$  in (3.4), that is,  $\varphi(t) = \phi_{n-1}(t)$  in (3.5), namely,  $d_k = d_k^{(0)}$  ( $0 \leq k \leq$

$n - 1$ ). We set  $d_k = 0$  ( $k \geq n$ ). Forward and backward computations of (1.13) are numerically stable for  $0 \leq k \leq n - 1$ . With the conditions  $d_n = 0$  and  $d_{n-1} = b_{n-1}$ , backward recursion of (1.13) gives  $d_{n-2}, \dots, d_0$ . Then, we have

$$I(p_n; c, \omega) = \frac{e^{i\omega}}{i\omega} \sum_{k=0}^{n-1} {}' d_k - \frac{e^{-i\omega}}{i\omega} \sum_{k=0}^{n-1} {}' (-1)^k d_k. \quad (3.8)$$

When  $n - 1 > M$ , neither forward nor backward computation of (1.13) is stable for  $M < k \leq n - 1$ . In the computation process the value of  $d_k$  is soon contaminated with the rounding error; see [8, 9]. In this case, we need a solution of the second type ( $C \neq 0$ ). We use the infinite Chebyshev series  $\varphi(x)$  given by  
125 (1.11), where the coefficients  $d_k$  satisfy the RR (1.13). To avoid the numerical instability, we reformulate the RR (1.13) as the problem of a system of linear equations. The normalization relation  $\varphi(-1) = \sum_{k=0}^{\infty} {}' (-1)^k d_k = 0$  is used in [8].

In this paper, rather than this relation, we prefer the condition  $d_M = 0$  (cf.  
130 [7]) that leads to reduction of the computational cost; see Section 4. Then, from (3.7) we obtain  $C = -d_M^{(0)} / \{2(-i)^M J_M(\omega)\}$ . We assure that  $|C| < \infty$ , since  $|d_M^{(0)}| < \infty$  and  $J_M(\omega) \neq 0$ , see Appendix B.

Since  $d_k$  converges rapidly to 0 ( $k \rightarrow \infty$ ) as shown above, instead of the infinite summation (1.11) for  $\varphi(t)$ , we use the truncated Chebyshev series (1.15). We determine  $d_k^{[N]}$  ( $0 \leq k \leq N$ ) so that (1.13) is satisfied for  $1 \leq k \leq N$  with the condition  $d_M^{[N]} = d_M = 0$  and  $d_k^{[N]} = 0$  ( $k > N$ ), namely,

$$d_{k-1}^{[N]} + \frac{2k}{i\omega} d_k^{[N]} - d_{k+1}^{[N]} = b_{k-1} - b_{k+1} \quad (1 \leq k \leq N), \quad (3.9)$$

recalling that  $b_k = 0$  ( $k \geq n$ ). Then, we have the approximation  $I^{[N]}(p_n; c, \omega)$  (1.16). Section 4 shows a system of linear equations converted from (3.9) for  
135  $M + 1 \leq k \leq N$ . We solve the linear system by Gaussian elimination with the smallest  $N$  satisfying the condition  $|d_N^{[N]}| \leq u$ , namely,  $E_T \leq u$  from Lemma 1.1. Then, with  $d_{M+1}^{[N]}$  obtained above and the condition  $d_M^{[N]} = 0$ , backward computation of the RR (3.9) for  $1 \leq k \leq M$  gives  $d_k^{[N]}$  ( $0 \leq k \leq M - 1$ ).

#### 4. Computing $d_k^{[N]}$

140 Recall that  $M = \lfloor \omega \rfloor$ . The linear system converted from the RR (3.9) for  $M+1 \leq k \leq N$  with the condition  $d_M^{[N]} = 0$  is solved in the same way as the one given in [7]. So, we omit the details and show only a difference (the stopping criterion) and the computational costs required to compute  $d_k^{[N]}$  ( $0 \leq k \leq N$ ).

According to the notation in [7], let  $\mu_k = 2k/\omega$  and  $\alpha_k = b_{k-1} - b_{k+1}$  ( $1 \leq k \leq N$ ). Letting  $d_M^{[N]} = 0$ , we rewrite (3.9) as follows,

$$d_{k-1}^{[N]} - i\mu_k d_k^{[N]} - d_{k+1}^{[N]} = \alpha_k \quad (1 \leq k \leq M), \quad (4.1)$$

$$\left. \begin{aligned} -i\mu_{M+1} d_{M+1}^{[N]} - d_{M+2} &= \alpha_{M+1}, \\ d_{k-1}^{[N]} - i\mu_k d_k^{[N]} - d_{k+1}^{[N]} &= \alpha_k \quad (M+2 \leq k \leq N-1), \\ d_{N-1}^{[N]} - i\mu_N d_N^{[N]} &= \alpha_N. \end{aligned} \right\} \quad (4.2)$$

Note that  $\mu_k$  ( $1 \leq k \leq N$ ) are stored in a real array.

145 We solve the tridiagonal linear system (4.2) by Gaussian elimination without pivoting, increasing  $N$  until the stopping criterion  $|d_N^{[N]}| \leq u$  is satisfied in the forward elimination process. This means that  $E_T \leq u$ , from Lemma 1.1.

Keller [7] adopts a different stopping criterion. If we modify his stopping rule in Algorithm 2.4 in [7] to  $|\tilde{\alpha}_{j+1}/\tilde{\mu}_{j+1}| \leq u$  for some  $j \geq M$  ( $\tilde{\alpha}_{j+1}$  and  $\tilde{\mu}_{j+1}$  150 defined in [7]) and set  $N = j+1$ , then we obtain practically the same algorithm as ours<sup>1</sup>. The value of  $d_N^{[N]}$ , that is,  $-\tilde{\alpha}_N/(i\tilde{\mu}_N)$ , is obtained in the forward elimination (2.11) in [7], so requires no additional computation.

Finally, we compute the RR (4.1) in the backward direction to obtain  $d_k^{[N]}$  ( $0 \leq k \leq M-1$ ).

155 This process in [7, Lemma 2.2 and Lemma 2.3] to solve (4.2) followed by (4.1) requires a small number of arithmetic operations, namely,  $2N - 2M - 1$  complex/reals,  $M - 1$  complex $\times$ reals and  $N - M - 1$  real reciprocals.

Our previous scheme in [8] requires  $11N - 8M - 3$  complex multiplication-divisions.

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<sup>1</sup>The authors thank the reviewer for pointing this out.

160 **5. Algorithm**

We give the algorithm to compute  $Q(p_n; c, \omega)$  in (1.6). We assume the double precision computation with the unit roundoff  $u \cong 1.11 \times 10^{-16}$ . We assume that standard routines for the FFT,  $\text{Ci}(x)$  and  $\text{Si}(x)$  are available.

**Purpose**

165 Given the function  $f$ , the values of  $c \in [-1, 1]$  and  $\omega$  and the number of function evaluations  $n + 1 \geq 2$  required, we compute an approximation  $Q(p_n; c, \omega)$ .

**Outline of the algorithm**

**Output** is the approximation  $Q(p_n; c, \omega)$ .

Step 1. Initialization. Set  $u = 2^{-53} \cong 1.11 \times 10^{-16}$ .

170 Step 2. Use the FFT<sup>†</sup> to compute the summations on the right-hand side of (3.1) to obtain the Chebyshev coefficients  $a_k$  ( $0 \leq k \leq n$ ) of  $p_n$ .

Step 3. Compute the RR (3.2) in the backward direction with starting values  $b_n = 0$  and  $b_{n-1} = a_n/2$  to obtain  $b_k$  ( $0 \leq k \leq n - 1$ ).

Step 4. If  $n - 1 \leq |\omega|$ , go to Step 5. Otherwise, go to Step 6.

175 Step 5. Compute the RR (1.13) in the backward direction with starting values  $d_n = d_{n+1} = 0$  to obtain  $d_k$  ( $0 \leq k \leq n - 1$ ).  
Compute the right-hand side of (3.8) to obtain  $I(p_n; c, \omega)$ .  
Go to Step 8.

Step 6. Solve the linear system (4.2) for  $d_k^{[N]}$  by Gaussian elimination via bordering with the smallest  $N$  that satisfies  $|d_N^{[N]}| \leq u$ . Set  $d_M^{[N]} = 0$ . Compute  
180 the RR (4.1) in the backward direction to obtain  $d_k^{[N]}$  ( $0 \leq k \leq M - 1$ ).

Step 7. Compute the approximation  $I^{[N]}(p_n; c, \omega)$  in (1.16) to  $I(p_n; c, \omega)$ .

Step 8. If  $c \in (-1, 1)$ , compute  $\text{Ci}(|(1 - c)\omega|)$ ,  $\text{Ci}(|(1 + c)\omega|)$ ,  $\text{Si}((1 - c)\omega)$  and  
185  $\text{Si}((1 + c)\omega)$  to obtain  $\int_{-1}^1 e^{ic\omega} dt / (t - c)$  in (2.1). If  $c = \pm 1$ , compute  
 $\text{Ci}(2|\omega|)$  and  $\text{Si}(2\omega)$  to obtain  $\int_{-1}^1 e^{i\omega t} / (t - c) dt$  in (2.3).

Step 9. Evaluate  $p_n(c)$  by Clenshaw's algorithm<sup>‡</sup>.

Step 10. Evaluate  $Q(p_n; c, \omega)$  by (1.6), where  $I(p_n; c, \omega)$  is replaced by the approximation  $I^{[N]}(p_n; c, \omega)$  in Step 7 when  $n - 1 > |\omega|$ .

† The *discrete cosine transform* (DCT type 1) with the FFT (cf. [19, p.229, 190 p.232]) for  $a_k = a(k+1)$  ( $0 \leq k \leq n$ ), Chebyshev coefficients of the polynomial  $p_n$  interpolating  $f$ , is written as a Matlab script,

```
t=linspace(0,pi,n+1); t=cos(t); y=f(t);
g=[y y(n:-1:2)]; g=fft(g')/n; a=g(1:n+1);
```

‡ *Clenshaw's algorithm* (cf. [24, p.166],[3, p.27],[34, p.157]) for evaluating  $p_n(c)$  at  $t = c$  in (1.2) is as follows. Set  $y_{n+1} = 0$  and  $y_n = a_n/2$ . Compute  $y_k$  by

$$y_k = 2c y_{k+1} - y_{k+2} + a_k \quad (k = n-1, \dots, 0). \quad (5.1)$$

Then,  $p_n(c) = (y_0 - y_2)/2$ .

## 195 6. Error analysis

We estimate the error of the approximation  $Q(p_n; c, \omega)$  in (1.6) to the integrals  $Q(f; c, \omega)$  in (1.1). For  $f \in C^1[-1, 1]$ , the uniform convergence of the method is proved in [13]. Here, assuming that  $f$  is analytic on a region in the complex plane, we show the uniform convergence using a tighter error bound. A 200 similar error analysis is in [12], see also [11]. See [35] for uniform approximations to hypersingular finite-part integrals.

### 6.1. Uniform error bound

Let  $\mathcal{E}_\rho$  denote an ellipse in the complex plane:  $|z-1| + |z+1| = \rho + \rho^{-1}$ , whose foci are at  $z = \pm 1$  and the sum of semi-axes is  $\rho > 1$ . Let  $u = \rho e^{i\theta}$  ( $0 \leq \theta \leq 2\pi$ ), then we have

$$\mathcal{E}_\rho : z = \frac{1}{2}(u + u^{-1}). \quad (6.1)$$

Let  $W_{n+1}(x)$  be a polynomial of degree  $n+1$  defined by

$$W_{n+1}(t) = T_{n+1}(t) - T_{n-1}(t) = 2(t^2 - 1)U_{n-1}(t) \quad (n \geq 1). \quad (6.2)$$

Then, the Chebyshev points  $\cos(\pi j/n)$  ( $0 \leq j \leq n$ ) are zeros of  $W_{n+1}(x)$ . Lemma 6.1 is given in [3, p.150].

**Lemma 6.1** (cf. [3]). *Suppose that  $f(z)$  is single-valued and analytic inside and on  $\mathcal{E}_\rho$  defined by (6.1). Then, the error  $e_n(t)$  of the interpolating polynomial  $p_n(t)$  in (1.2) is expressed in terms of a contour integral,*

$$e_n(t) = f(t) - p_n(t) = \frac{W_{n+1}(t)}{2\pi i} \oint_{\mathcal{E}_\rho} \frac{f(z)}{W_{n+1}(z)(z-t)} dz. \quad (6.3)$$

205 **Remark 6.2.** *If  $f(z)$  is analytic on  $[-1, 1]$ , then there exists a sufficiently small  $\rho > 1$  such that  $f(z)$  is analytic inside and on the ellipse  $\mathcal{E}_\rho$ .*

Let  $E_1$  and  $E_2$  be defined by

$$E_1 = \int_{-1}^1 \frac{e_n(t) - e_n(c)}{t-c} e^{i\omega t} dt, \quad (6.4)$$

$$E_2 = \begin{cases} e_n(c) \int_{-1}^1 \frac{e^{i\omega t}}{t-c} dt, & c \in (-1, 1), \\ 0, & c = \pm 1, \end{cases} \quad (6.5)$$

respectively. Then, the error of the approximation  $Q(p_n; c, \omega)$  is given by

$$|Q(f; c, \omega) - Q(p_n; c, \omega)| = |Q(e_n; c, \omega)| = |E_1 + E_2| \leq |E_1| + |E_2|. \quad (6.6)$$

Now, let us bound  $E_1$  and  $E_2$ .

**Lemma 6.3.** *Suppose that  $f(z)$  is single-valued and analytic inside and on  $\mathcal{E}_\rho$  defined by (6.1) and let  $M = \max_{z \in \mathcal{E}_\rho} |f(z)|$ . Then, for  $c \in [-1, 1]$ , we bound  $E_1$  in (6.4) by*

$$|E_1| \leq \frac{8M\rho}{\pi(\rho-1)^2(\rho^n - \rho^{-n})} \left\{ 2(2n+1) + \pi \log \left( \frac{\rho+1}{\rho-1} \right) \right\} \quad (c \in [-1, 1]). \quad (6.7)$$

*Proof.* In view of Lemma 6.1 and (6.4) we obtain

$$E_1 = \frac{1}{2\pi i} \oint_{\mathcal{E}_\rho} \frac{f(z)}{W_{n+1}(z)} \int_{-1}^1 \frac{1}{t-c} \left\{ \frac{W_{n+1}(t)}{z-t} - \frac{W_{n+1}(c)}{z-c} \right\} e^{i\omega t} dt dz. \quad (6.8)$$

Let  $G_1$  be defined by

$$G_1 = \max_{\substack{z \in \mathcal{E}_\rho \\ c \in [-1, 1]}} \int_{-1}^1 \left| \frac{1}{t-c} \left\{ \frac{W_{n+1}(t)}{z-t} - \frac{W_{n+1}(c)}{z-c} \right\} \right| dt. \quad (6.9)$$

Then, from (6.8) we have

$$|E_1| \leq \frac{MG_1}{2\pi} \oint_{\mathcal{E}_\rho} \left| \frac{dz}{W_{n+1}(z)} \right|. \quad (6.10)$$



We claim that

$$\oint_{\mathcal{E}_\rho} \left| \frac{dz}{W_{n+1}(z)} \right| \leq \frac{2\pi}{\rho^n - \rho^{-n}}, \quad (6.11)$$

and

$$G_1 \leq \frac{8\rho}{\pi(\rho-1)^2} \left\{ 2(2n+1) + \pi \log \left( \frac{\rho+1}{\rho-1} \right) \right\}. \quad (6.12)$$

Then, inserting (6.11) and (6.12) into (6.10) proves the lemma.

It remains to verify (6.11) and (6.12). Verification of (6.11) is simple. For

$$z = \frac{1}{2}(u + u^{-1}) \in \mathcal{E}_\rho, \quad u = \rho e^{i\theta},$$

(6.1), we have  $T_n(z) = \frac{1}{2}(u^n + u^{-n})$  (cf. [3, p.14]). Consequently, from (6.2) we have

$$W_{n+1}(z) = \frac{1}{2}(u^{n+1} + u^{-n-1}) - \frac{1}{2}(u^{n-1} + u^{-n+1}) = \frac{1}{2}(u^n - u^{-n})(u - u^{-1}). \quad (6.13)$$

We verify (6.11) since in view of  $dz = \frac{1}{2}(1 - u^{-2})du$  and (6.13) we have

$$\oint_{\mathcal{E}_\rho} \left| \frac{dz}{W_{n+1}(z)} \right| = \oint_{|u|=\rho} \left| \frac{du}{(u^n - u^{-n})u} \right| \leq \oint_{|u|=\rho} \frac{du}{(\rho^n - \rho^{-n})\rho} = \frac{2\pi}{\rho^n - \rho^{-n}}. \quad (6.14)$$

Verification of (6.12) requires some manipulations. Since

$$\frac{1}{t-c} \left\{ \frac{W_{n+1}(t)}{z-t} - \frac{W_{n+1}(c)}{z-c} \right\} = \frac{W_{n+1}(t) - W_{n+1}(c)}{(z-t)(t-c)} + \frac{W_{n+1}(c)}{(z-t)(z-c)},$$

$$\min_{\substack{z \in \mathcal{E}_\rho \\ t \in [-1,1]}} |z-t| = \frac{\rho + \rho^{-1}}{2} - 1 = \frac{(\rho-1)^2}{2\rho}, \quad (6.15)$$

$$\max_{z \in \mathcal{E}_\rho} \int_{-1}^1 \frac{dt}{|z-t|} = 2 \log \left( \frac{\rho+1}{\rho-1} \right), \quad (6.16)$$

see (7.12) in [8] for (6.16), and  $|W_{n+1}(c)| \leq 2$  ( $c \in [-1, 1]$ ), from (6.9) we have

$$\begin{aligned} G_1 &\leq \max_{\substack{z \in \mathcal{E}_\rho \\ c \in [-1,1]}} \int_{-1}^1 \left| \frac{W_{n+1}(t) - W_{n+1}(c)}{(z-t)(t-c)} \right| dt + \max_{\substack{z \in \mathcal{E}_\rho \\ c \in [-1,1]}} \left| \frac{W_{n+1}(c)}{z-c} \right| \int_{-1}^1 \frac{dt}{|z-t|} \\ &\leq \frac{2\rho}{(\rho-1)^2} \left\{ \max_{c \in [-1,1]} \int_{-1}^1 \left| \frac{W_{n+1}(t) - W_{n+1}(c)}{t-c} \right| dt + 4 \log \left( \frac{\rho+1}{\rho-1} \right) \right\}. \end{aligned} \quad (6.17)$$

Since

$$U_k(t) - U_{k-2}(t) = 2T_k(t) \quad (k \geq 2), \quad U_1(t) = 2T_1(t),$$

(cf. [34, p.9]), and

$$T_{n+1}(t) - T_{n+1}(c) = 2(t-c) \sum_{k=0}^n U_{n-k}(t) T_k(c),$$

(cf. [36]), from (6.2) we have

$$\frac{W_{n+1}(t) - W_{n+1}(c)}{t-c} = 4 \sum_{j=0}^n T_{n-j}(c) T_j(t). \quad (6.18)$$

In view of (6.18) and the relations

$$\int_{-1}^1 |T_0(t)| dt = 2 < \frac{8}{\pi}, \quad \int_{-1}^1 |T_1(t)| dt = 1 < \frac{4}{\pi}, \quad \int_{-1}^1 |T_k(t)| dt \leq \frac{4}{\pi} \quad (k \geq 2),$$

(cf. [34, p.38]), we have

$$\int_{-1}^1 \left| \frac{W_{n+1}(t) - W_{n+1}(c)}{t-c} \right| dt \leq 4 \sum_{j=0}^n \int_{-1}^1 |T_j(t)| dt \leq \frac{8(2n+1)}{\pi}. \quad (6.19)$$

Using (6.19) in (6.17) verifies (6.12).  $\square$

**Lemma 6.4.** *Under the same assumption as in Lemma 6.3, we bound  $E_2$  in (6.5) by*

$$|E_2| \leq \frac{16M\rho}{(\rho-1)^2(\rho^n - \rho^{-n})} \quad (c \in [-1, 1]). \quad (6.20)$$

*Proof.* When  $c = \pm 1$ , in view of  $W_{n+1}(c) = 0$  we have  $e_n(c) = 0$  in (6.3), consequently  $E_2 = 0$  in (6.5). Now, we assume that  $c \in (-1, 1)$ . From Lemma 2.1 and (6.5) we have

$$|E_2| \leq |e_n(c)| \times \sqrt{\{\text{Ci}(|(1-c)\omega|) - \text{Ci}(|(1+c)\omega|)\}^2 + \{\text{Si}((1-c)\omega) + \text{Si}((1+c)\omega)\}^2}. \quad (6.21)$$

In view of (6.3), (6.14) and (6.15) we have

$$\begin{aligned} |e_n(c)| &\leq \frac{|W_{n+1}(c)|}{2\pi} \oint_{\mathcal{E}_\rho} \left| \frac{f(z) dz}{W_{n+1}(z)(z-c)} \right| \\ &\leq \frac{M\rho |W_{n+1}(c)|}{\pi(\rho-1)^2} \oint_{\mathcal{E}_\rho} \left| \frac{dz}{W_{n+1}(z)} \right| \leq \frac{2M\rho |W_{n+1}(c)|}{(\rho-1)^2(\rho^n - \rho^{-n})}. \end{aligned} \quad (6.22)$$

In view of Lemma 2.1 we see that

$$\begin{aligned} |\operatorname{Ci}(|(1-c)\omega|) - \operatorname{Ci}(|(1+c)\omega|)| &= \left| \int_{|(1+c)\omega|}^{|(1-c)\omega|} \frac{\cos t}{t} dt \right| \\ &\leq \left| \int_{|(1+c)\omega|}^{|(1-c)\omega|} \frac{1}{t} dt \right| = \left| \log \frac{1-c}{1+c} \right| \quad (c \in (-1, 1)). \end{aligned} \quad (6.23)$$

In [32, p.244] it is shown that  $0 \leq \operatorname{Si}(t) \leq \operatorname{Si}(\pi) = 1.8519370\dots$ . From (6.2) setting  $c = \cos \vartheta$  we have

$$|W_{n+1}(c)| = 2|\sin \vartheta \sin n\vartheta| \leq 2|\sin \vartheta| = 2\sqrt{1-c^2}.$$

Using these results and (6.22) and (6.23) in (6.21) we have

$$|E_2| \leq \frac{4M\rho}{(\rho-1)^2(\rho^n - \rho^{-n})} \sqrt{\phi(c)}, \quad (6.24)$$

$$\phi(c) = (1-c^2) \left\{ \left( \log \frac{1-c}{1+c} \right)^2 + 4s^2 \right\}, \quad (6.25)$$

210 where we set  $s = \operatorname{Si}(\pi)$ . We claim that  $\max_{-1 < c < 1} \phi(c) = 4s^2 < 16$ . Then, (6.20) follows from (6.24).

Verifying that  $\max_{-1 < c < 1} \phi(c) = 4s^2$  is simple. Since  $\phi(c)$  is an even function, it is enough to examine only the case,  $0 \leq c < 1$ . Since  $\phi(0) = 4s^2$ ,  $\phi'(0) = 0$ ,  $\phi''(0) = 8(1-s^2) < 0$  and

$$\phi^{(3)}(c) = \frac{16 \log\{(1-c)/(1+c)\}}{(1-c^2)^2} < 0 \quad (0 < c < 1),$$

we have  $\max_{0 \leq c \leq 1} \phi(c) = \phi(0) = 4s^2$ . □

In view of (6.6), Lemmas 6.3 and 6.4, Theorem 6.5 for the error bound of the approximation  $Q(p_n; c, \omega)$  immediately follows.

**Theorem 6.5.** *Under the same assumption as in Lemma 6.3, the error of the approximations  $Q(p_n; c, \omega)$  is bounded uniformly by*

$$\begin{aligned} |Q(f; c, \omega) - Q(p_n; c, \omega)| &\leq \frac{8M\rho [2(2n+1+\pi) + \pi \log\{(\rho+1)/(\rho-1)\}]}{\pi(\rho-1)^2(\rho^n - \rho^{-n})} \\ &\sim (32M/\pi) n\rho^{-n-1} = O(n\rho^{-n}) \quad (n \rightarrow \infty). \end{aligned} \quad (6.26)$$

215 While our error bound does not depend on  $\omega$ , in the approximation method given in [17] that is of a different type from ours, the absolute error of the approximation decreases very fast as  $|\omega|$  grows.

6.2. An illustrative evidence of uniform error bound

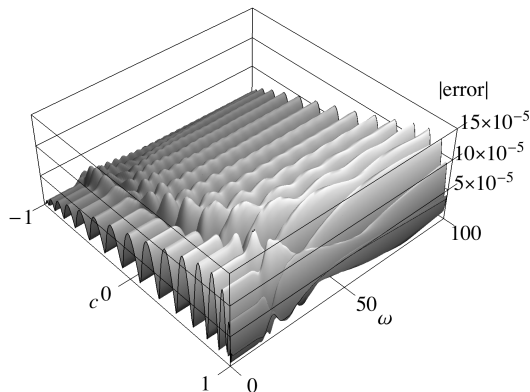


Figure 1: The error of the approximation at  $n + 1 = 17$  nodes for the integral  $Q(f; c, \omega)$  as a function of  $c$  and  $\omega$ , where  $f(t) = (1 - \alpha^2)/(1 - 2\alpha t + \alpha^2)$  and  $\alpha = 0.5$ .

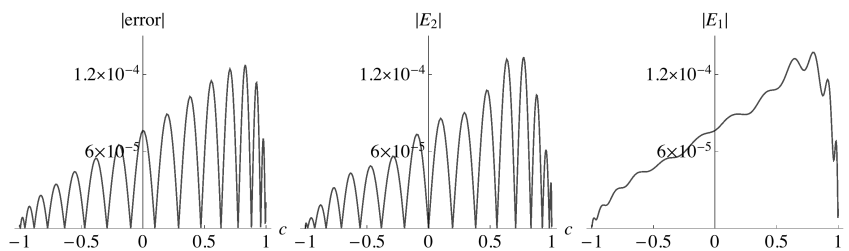


Figure 2: The error (*left*) of the approximation at  $n + 1 = 17$  nodes for the integral  $Q(f; c, 10)$  as a function of  $c$ ,  $|E_2|$  (6.5) (*center*) and  $|E_1|$  (6.4) (*right*), where  $f(t) = (1 - \alpha^2)/(1 - 2\alpha t + \alpha^2)$  and  $\alpha = 0.5$ .

Figure 1 illustrates the absolute error  $|Q(f; c, \omega) - Q(p_n; c, \omega)|$  of the approximation  $Q(p_n; c, \omega)$  at  $n + 1 = 17$  nodes to the integral  $Q(f; c, \omega)$  for  $f(t) = (1 - \alpha^2)/(1 - 2\alpha t + \alpha^2)$  with  $\alpha = 0.5$ ,  $c \in [-1, 1]$  and  $1 \leq \omega \leq 100$ . Figure 2 shows the error  $|Q(f; c, \omega) - Q(p_n; c, \omega)| = |Q(e_n; c, \omega)|$ ,  $|E_2|$  and  $|E_1|$  for  $\omega = 10$  and Figure 3 for  $\omega = 100$ . (Recall that the quadrature error in magnitude is  $|Q(e_n; c, \omega)| = |E_1 + E_2| \leq |E_1| + |E_2|$ ). We observe that the error is uniformly bounded, independently of the values of  $c$  and  $\omega$ . Further,

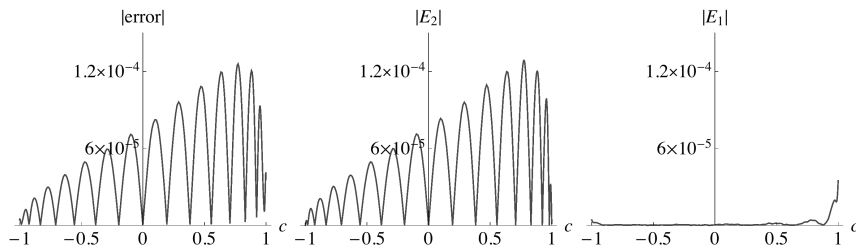


Figure 3: The error (left) of the approximation at  $n+1 = 17$  nodes for the integral  $Q(f; c, 100)$  as a function of  $c$ ,  $|E_2|$  (6.5) (center) and  $|E_1|$  (6.4) (right), where  $f(t) = (1-\alpha^2)/(1-2\alpha t+\alpha^2)$  and  $\alpha = 0.5$ .

$|E_2|$  vanishes when  $c$  coincides with the nodes, where  $e_n(c) = 0$  in (6.5). From Figures 2 and 3 we observe that  $|E_1|$  and  $|E_2|$  are comparable for small  $|\omega|$  while  $|E_1|$  is much smaller than  $|E_2|$  for larger  $|\omega|$ , indeed,  $|E_1| = O(1/\omega)$  as  $|\omega| \rightarrow \infty$ . Consequently,  $|E_2|$  dominantly contributes to the quadrature error (6.6), particularly for large  $|\omega|$ .

## 7. Numerical examples

In this section, by numerical examples we demonstrate the performance of our algorithm. The numerical computations are carried out with about 15 significant digits by using Mathematica (Ver. 11.1.1.0) on the iMac (3.2GHz, Intel Core i5). The values that we regard as exact are obtained with 100 significant digits computation on Mathematica.

We compute five types of oscillatory CPV integrals,

$$Q(f_m; c, \omega) = \int_{-1}^1 \frac{f_m(t) e^{i\omega t}}{t - c} dt, \quad m = 1, \dots, 5,$$

for  $c = 0.9$  and  $\omega = 10, 1000$ , where the functions  $f_m(t)$  are given by

$$f_1(t) = e^{\alpha(t-1)}, \quad \alpha = 4, 16, \quad (7.1)$$

$$f_2(t) = e^{i2\pi\alpha t}, \quad \alpha = 8, 16, \quad (7.2)$$

$$f_3(t) = \frac{1 - \alpha^2}{1 - 2\alpha t + \alpha^2}, \quad \alpha = 0.8, 0.9, \quad (7.3)$$

$$f_4(t) = \frac{1}{t^2 + \alpha^2}, \quad \alpha = 1/4, 1/8, \quad (7.4)$$

$$f_5(t) = (1 - t^2)^{3/2}, \quad (7.5)$$

respectively. We compute oscillatory HFP integrals with  $c = 1$  and  $\omega = 10, 1000$  as well and omit the details since obtained results for relative errors are very similar to Figures 4~8 below for the oscillatory CPV integrals. We note that  
 240 the uniform dependence of the approximation  $Q(p_n; c, \omega)$  on the value of  $n$  is common for most Clenshaw-Curtis like methods, including [12–14].

The function in (7.1) is an entire function, (7.2) an oscillatory function, (7.3) the generating function of the Chebyshev polynomial (cf. [22], [34, p.41]), (7.4) a  
 245 peaked function and (7.5) a function having singularities of the second-derivative at both ends of the integration interval.

The integrals  $Q(f_m; 0.9, \omega)$  ( $m = 1, \dots, 4$ ) except for  $Q(f_5; 0.9, \omega)$  with their exact values (for  $\omega = 10, 1000$ ) are expressed by

$$\begin{aligned} Q(f_1; 0.9, 10) &= \Phi_I(0.9, 10 - i\alpha) e^{-\alpha} \\ &= \begin{cases} -1.1256339442498735738 - 1.2174807464660865793i & (\alpha = 4), \\ -0.79432599720832534426 - 0.27066674690448630758i & (\alpha = 16), \end{cases} \end{aligned}$$

$$\begin{aligned} Q(f_1; 0.9, 1000) &= \Phi_I(0.9, 1000 - i\alpha) e^{-\alpha} \\ &= \begin{cases} -2.0930127016937171914 + 0.1338344041044348488i & (\alpha = 4), \\ -0.62458187202084588291 + 0.03644532053890920537i & (\alpha = 16), \end{cases} \end{aligned}$$

$$\begin{aligned} Q(f_2; 0.9, 10) &= \Phi_R(0.9, 10 + 2\pi\alpha) \\ &= \begin{cases} 2.2610032378003899837 - 1.9621437455729568795i & (\alpha = 8), \\ 2.6905839749162102713 + 1.6377598933795257917i & (\alpha = 16), \end{cases} \end{aligned}$$

$$\begin{aligned}
Q(f_2; 0.9, 1000) &= \Phi_R(0.9, 1000 + 2\pi\alpha) \\
&= \begin{cases} -1.1592005127547868638 - 2.9226666731058770971i & (\alpha = 8), \\ 2.4207591820350711974 - 2.0163447425643565369i & (\alpha = 16), \end{cases}
\end{aligned}$$

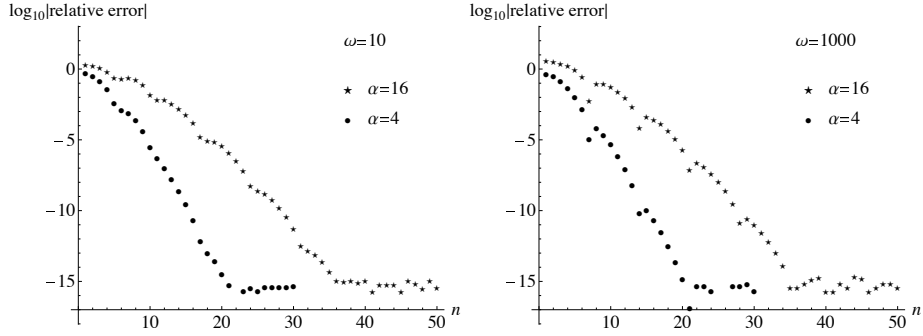


Figure 4: The relative errors of the approximations to the integrals  $Q(f_1; 0.9, \omega)$  for  $f_1(t) = e^{\alpha(t-1)}$  with  $\alpha = 4, 16$  and for  $\omega = 10$  (left) and 1000 (right).

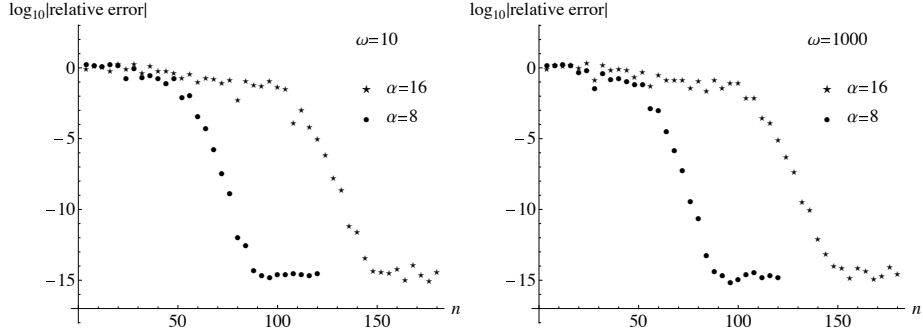


Figure 5: The relative errors of the approximations to the integrals  $Q(f_2; 0.9, \omega)$  for  $f_2(t) = e^{i2\pi\alpha t}$  with  $\alpha = 8, 16$  and for  $\omega = 10$  (left) and 1000 (right).

$$\begin{aligned}
Q(f_3; 0.9, 10) &= \frac{1 - \alpha^2}{2\alpha(r - c)} \{ \Phi_C(r, 10) - \Phi_R(0.9, 10) \} \\
&= \begin{cases} -5.1726063581663838102 - 3.3490166983645622408i & (\alpha = 0.8), \\ -4.1057284162415934661 - 2.2927376525987158576i & (\alpha = 0.9), \end{cases} \\
r &:= (\alpha + 1/\alpha)/2,
\end{aligned}$$

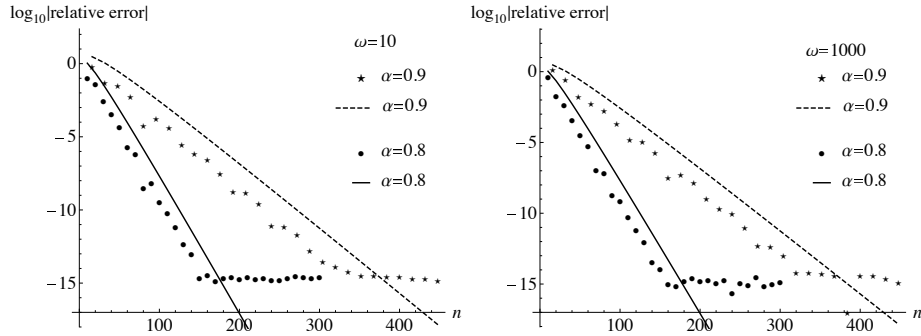


Figure 6: The relative errors of the approximations to the integrals  $Q(f_3; 0.9, \omega)$  for  $f_3(t) = (1 - \alpha^2)/(1 - 2\alpha t + \alpha^2)$  with  $\alpha = 0.8, 0.9$  and for  $\omega = 10$  (*left*) and  $1000$  (*right*). The solid and broken curves represent  $n\rho^{-n}$  for the parameter  $\rho$  of the ellipse  $\mathcal{E}_\rho$  (6.1).

$$\begin{aligned}
 Q(f_3; 0.9, 1000) &= \frac{1 - \alpha^2}{2\alpha(r - c)} \{ \Phi_C(r, 1000) - \Phi_R(0.9, 1000) \} \\
 &= \begin{cases} -5.5667603364747801020 + 0.3263089276087251178i & (\alpha = 0.8), \\ -2.9692938030411933993 + 0.1295150056147923552i & (\alpha = 0.9), \end{cases} \\
 r &:= (\alpha + 1/\alpha)/2,
 \end{aligned}$$

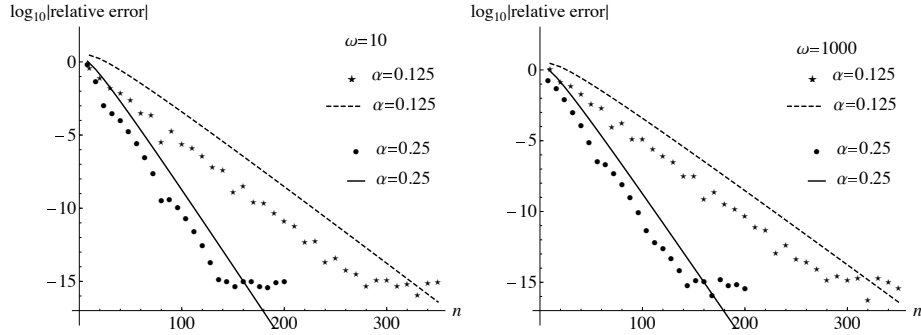


Figure 7: The relative errors of the approximations to the integrals  $Q(f_4; 0.9, \omega)$  for  $f_4(t) = 1/(t^2 + \alpha^2)$  with  $\alpha = 1/4, 1/8$  and for  $\omega = 10$  (*left*) and  $1000$  (*right*). The solid and broken curves represent  $n\rho^{-n}$  for the parameter  $\rho$  of the ellipse  $\mathcal{E}_\rho$  (6.1).



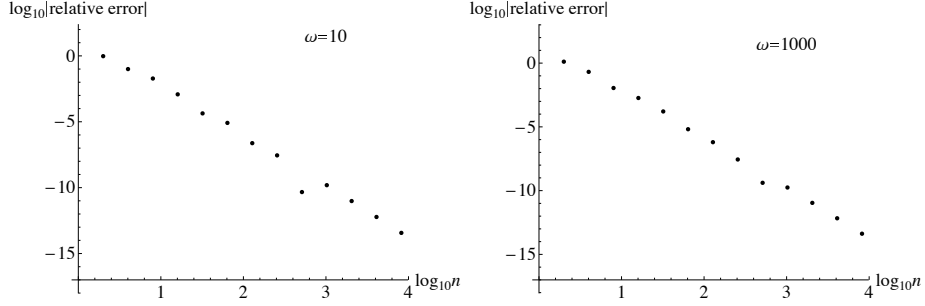


Figure 8: The relative errors of the approximations to the integrals  $Q(f_5; 0.9, \omega)$  for  $f_5(t) = (1 - t^2)^{3/2}$  and for  $\omega = 10$  (left) and 1000 (right).

$$\begin{aligned}
 Q(f_4; 0.9, 10) &= \frac{1}{2i\alpha} \left\{ \frac{\Phi_C(i\alpha, 10) - \Phi_R(0.9, 10)}{i\alpha - 0.9} + \frac{\Phi_C(-i\alpha, 10) - \Phi_R(0.9, 10)}{i\alpha + 0.9} \right\} \\
 &= \begin{cases} -2.5024802215496231205 - 2.9010231004723840715i & (\alpha = 1/4), \\ -9.3674757276956166055 - 3.8530400190128827327i & (\alpha = 1/8), \end{cases}
 \end{aligned}$$

$$\begin{aligned}
 Q(f_4; 0.9, 1000) &= \frac{1}{2i\alpha} \left\{ \frac{\Phi_C(i\alpha, 1000) - \Phi_R(0.9, 1000)}{i\alpha - 0.9} \right. \\
 &\quad \left. + \frac{\Phi_C(-i\alpha, 1000) - \Phi_R(0.9, 1000)}{i\alpha + 0.9} \right\} \\
 &= \begin{cases} -3.5854608269985750047 + 0.2328694521987087191i & (\alpha = 1/4), \\ -3.7891040784328703465 + 0.2461500798066019073i & (\alpha = 1/8), \end{cases}
 \end{aligned}$$

where  $\Phi_R(c, \omega)$ ,  $\Phi_I(c, \omega)$  and  $\Phi_C(c, \omega)$  are defined by

$$\begin{aligned}
 \Phi_R(c, \omega) &= \int_{-1}^1 \frac{e^{i\omega t}}{t - c} dt, \quad c \in (-1, 1), \omega \in \mathbb{R} \setminus \{0\}, \\
 \Phi_I(c, \omega) &= \int_{-1}^1 \frac{e^{i\omega t}}{t - c} dt, \quad c \in (-1, 1), \omega \in \mathbb{C} \setminus \mathbb{R}, \\
 \Phi_C(c, \omega) &= \int_{-1}^1 \frac{e^{i\omega t}}{t - c} dt, \quad c \in \mathbb{C} \setminus (-1, 1), \omega \in \mathbb{R},
 \end{aligned}$$

respectively and are written by

$$\begin{aligned}
 \Phi_R(c, \omega) &= e^{i\omega c} [\text{Ci}(|(1 - c)\omega|) - \text{Ci}(|(1 + c)\omega|) + iS] \\
 \Phi_I(c, \omega) &= e^{i\omega c} [\text{Ci}((1 - c)\omega) - \text{Ci}((1 + c)\omega) + iS], \\
 \Phi_C(c, \omega) &= e^{i\omega c} [\text{Ci}((1 - c)\omega) - \text{Ci}(-(1 + c)\omega) + iS],
 \end{aligned}$$

respectively. Here we set  $S = \text{Si}((1-c)\omega) + \text{Si}((1+c)\omega)$ . On the other hand, the exact value for the integral  $Q(f_5; c, \omega)$  is obtained as follows. First, we rewrite  $Q(f_5; c, \omega)$  as

$$Q(f_5; c, \omega) = \int_{-1}^1 \frac{f_5(t) - f_5(c)}{t - c} e^{i\omega t} dt + f_5(c) \Phi_R(c, \omega),$$

and compute the integral on the right-hand side of the equation above by using the routine `NIntegrate` in `Mathematica` with the required tolerance  $10^{-30}$ . Then, we have, for  $c = 0.9$  and  $\omega = 10, 1000$ ,

$$Q(f_5; 0.9, 10) = -0.08561094788020693315 - 0.28387269290696526616i$$

$$Q(f_5; 0.9, 1000) = -0.25961337137912856956 + 0.01723611261373691240i$$

From Figure 8, we observe that the error of the approximation to  $Q(f_5; c, \omega)$  of  $f_5(t)$  with singularities of the second-derivative decreases slowly like  $O(1/n^\alpha)$ ,  $\alpha \approx 4$  as  $n$  grows.

## 8. Concluding remarks

250 We presented an efficient quadrature rule of Clenshaw-Curtis type for approximating Cauchy principal value integrals of oscillatory functions and oscillatory Hadamard finite-part integrals. An algorithm was provided in a form easy to implement. We incorporated an improved version of the routine for the oscillatory integrals [8] into the algorithm. We proved that the error of  
 255 the approximation is bounded independently of the values of  $c \in [-1, 1]$  and  $\omega \neq 0$ . Numerical examples illustrated the efficiency of our method. The present method has an application to the approximations of oscillatory Hilbert transforms on the interval  $[0, \infty)$  and  $(-\infty, \infty)$  (cf. [14]).

## Appendix A. Proof of Lemma 1.1

We outline the proof of Lemma 1.1 that is almost the same as that in our previous method [8] with the normalization relation  $\varphi(-1) = 0$ . For  $\varphi(x)$  in

(1.11) and  $\varphi^{[N]}(x)$  in (1.15) let  $\varepsilon^{[N]}(x) = \varphi^{[N]}(x) - \varphi(x)$ . From the proof of Lemma 3.1 in [8] we derive

$$\varepsilon^{[N]}(1)e^{i\omega} - \varepsilon^{[N]}(-1)e^{-i\omega} = \frac{i\omega}{2}d_N^{[N]} \int_{-1}^1 U_N(x)e^{i\omega x} dx. \quad (\text{A.1})$$

Since in view of (1.12) and (1.16) we have

$$\begin{aligned} I(p_n; c, \omega) &= \frac{1}{i\omega} \{\varphi(1)e^{i\omega} - \varphi(-1)e^{-i\omega}\}, \\ I^{[N]}(p_n; c, \omega) &= \frac{1}{i\omega} \{\varphi^{[N]}(1)e^{i\omega} - \varphi^{[N]}(-1)e^{-i\omega}\}, \end{aligned}$$

from (A.1) it follows that

$$I(p_n; c, \omega) - I^{[N]}(p_n; c, \omega) = \frac{-1}{2}d_n^{[N]} \int_{-1}^1 U_N(x)e^{i\omega x} dx.$$

260 So, we verify Lemma 1.1 since  $|\int_{-1}^1 U_N(x)e^{i\omega x} dx| \leq 2$ , see [8].  $\square$

## Appendix B. Proving that $J_M(\omega) > 0$

For a given  $|x| > 0$ , let  $n = \lfloor |x| \rfloor$ . Then, for the Bessel function  $J_n(x)$  of order  $n$ , we prove that  $|J_n(x)| > 0$  for  $0 < |x| \leq n + 1$ . Since  $J_n(-x) = (-1)^n J_n(x)$  (see (B.2) below), it suffices to show that  $J_n(x) > 0$  for  $0 < x \leq n + 1$ . Since  $J_0(x) > 0$  ( $0 < x \leq 1$ ) (cf. [32, p.359]), we prove the case  $n \geq 1$ . We claim that for an arbitrary integer  $\nu \geq 1$ ,

$$J_\nu(x) > 0, \quad J'_\nu(x) > 0 \quad (0 < x \leq \nu), \quad (\text{B.1})$$

where  $J'_\nu(x)$  is the derivative of  $J_\nu(x)$ . Then, since  $J_n(x) = J_{n+2}(x) + 2J'_{n+1}(x)$  (cf. [32, 9.1.27]), we verify that  $J_n(x) > 0$  for  $0 < x \leq n + 1$ .

It remains to verify (B.1). Denote by  $j_{\nu,1}$  the smallest positive zero of  $J_\nu(x)$  and by  $j'_{\nu,1}$  that of  $J'_\nu(x)$ . Then,  $\nu \leq j'_{\nu,1} < j_{\nu,1}$  (cf. [32, 9.5.2]). Since (cf. [32, 9.1.10])

$$J_\nu(x) = \left(\frac{x}{2}\right)^\nu \sum_{k=0}^{\infty} \frac{(-x^2/2)^k}{k! \Gamma(\nu + k + 1)} = \frac{1}{\Gamma(\nu + 1)} \left(\frac{x}{2}\right)^\nu + O(x^{\nu+2}), \quad (\text{B.2})$$

for sufficiently small  $\epsilon > 0$ , we have  $J_\nu(\epsilon) > 0$ , consequently  $J_\nu(x) > 0$  for 265  $0 < x < j_{\nu,1}$ . In view of the fact that  $\nu < j_{\nu,1}$ , this verify the first relation of (B.1). Since from (B.2) we have  $J'_\nu(x) = (\nu/2)(x/2)^{\nu-1}/\Gamma(\nu + 1) + O(x^{\nu+1})$ , we similarly verify the second relation of (B.1).  $\square$

## Acknowledgements

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